

# Chapter 3

## Binary relations on a set

### 3.1 Equivalence relation

**Definition 3.1.** Let  $\mathcal{R}$  be a binary relation on  $E$ .  $\mathcal{R}$  is an equivalence relation if:

1.  $\mathcal{R}$  is *reflexive* :

$$\forall x \in E, x\mathcal{R}x$$

2.  $\mathcal{R}$  is *symmetric* :

$$\forall x, y \in E, x\mathcal{R}y \iff y\mathcal{R}x$$

3.  $\mathcal{R}$  is *transitive* :

$$\forall x, y, z \in E, [x\mathcal{R}y \wedge y\mathcal{R}z] \implies x\mathcal{R}z$$

**Example 3.1.** We consider the following relation on  $\mathbb{Z}$ :

$$\forall x, y \in \mathbb{Z}, x\mathcal{R}y \iff \exists k \in \mathbb{Z}, x - y = 2k$$

it is an equivalence relation .

1.  $\mathcal{R}$  is *reflexive* : Let  $x \in \mathbb{Z}$ , we have

$$x - x = 2 \times 0 \iff x\mathcal{R}x$$

Then,  $\mathcal{R}$  is reflexive.

2.  $\mathcal{R}$  is *symmetric* :

Let  $x, y \in \mathbb{Z}$ , we have

$$\begin{aligned} x\mathcal{R}y &\iff \exists k \in \mathbb{Z}, x - y = 2k \\ &\iff y - x = 2k' \quad (k' = -k \in \mathbb{Z}) \\ &\iff y\mathcal{R}x \end{aligned}$$

Then,  $\mathcal{R}$  is symmetric.

3.  $\mathcal{R}$  is *transitive* : Let  $x, y, z \in \mathbb{Z}$ , we have

$$\begin{aligned} x\mathcal{R}y \wedge y\mathcal{R}z &\iff \left\{ \begin{array}{l} \exists k \in \mathbb{Z}, x - y = 2k \dots\dots\dots(1) \\ \wedge \\ \exists k' \in \mathbb{Z}, y - z = 2k' \dots\dots\dots(2) \end{array} \right. \\ (1) + (2) &\implies x - z = 2k'', \quad (k'' = (k - k') \in \mathbb{Z}) \\ &\implies x\mathcal{R}z \end{aligned}$$

Then,  $\mathcal{R}$  is transitive.

So  $\mathcal{R}$  is an equivalence relation .

### 3.1.1 Equivalence class

**Definition 3.2.** If  $\mathcal{R}$  is an equivalence relation in a set  $E$ , the equivalence class of  $x \in E$  is the set

$$\dot{x} = \{y \in E / x\mathcal{R}y\}.$$

**notation 3.1.** We denote by  $E/\mathcal{R}$  ( the set of quotients of  $E$  by  $\mathcal{R}$ ) the set of equivalence classes of  $\mathcal{R}$

$$E/\mathcal{R} = \{\dot{x} / x \in E\}$$

**Example 3.2.** In the previous example, give  $\dot{x}$  and  $E/\mathcal{R}$

$$\begin{aligned}\dot{x} &= \{y \in \mathbb{Z} / x\mathcal{R}y\} \\ &= \{y \in \mathbb{Z} / x - y = 2k\} \\ &= \{x - 2k / k \in \mathbb{Z}\} \\ &= \{\dots, x - 4, x - 2, x, x + 2, x + 4, \dots\}\end{aligned}$$

$$\begin{aligned}\dot{0} &= \{y \in \mathbb{Z} / 0\mathcal{R}y\} \\ &= \{\dots, -4, -2, 0, 2, 4, \dots\}\end{aligned}$$

$$\begin{aligned}\dot{1} &= \{y \in \mathbb{Z} / 1\mathcal{R}y\} \\ &= \{\dots, -3, -1, 1, 3, \dots\}\end{aligned}$$

$$\dot{2} = \dot{0}$$

$$\mathbb{Z}/\mathcal{R} = \{\dot{x} / x \in \mathbb{Z}\}.$$

$$\mathbb{Z}/\mathcal{R} = \{\dot{0}, \dot{1}\}.$$

**Proposition 3.1.** Let  $\mathcal{R}$  be an equivalence relation in the set  $E$ . Then,

- $\forall x \in E, \dot{x} \subset E$ .
- $\forall x \in E, \dot{x} \neq \emptyset$ .
- $\forall x, y \in E, x\mathcal{R}y \implies \dot{x} = \dot{y}$ .

## 3.2 Order relation

**Definition 3.3.** Let  $\mathcal{R}$  be a binary relation on  $E$ . It 's an order relation if:

1.  $\mathcal{R}$  is *reflexive* :

$$\forall x \in E, x\mathcal{R}x.$$

2.  $\mathcal{R}$  is *anti symmetric* :

$$\forall x, y \in E, [x\mathcal{R}y \wedge y\mathcal{R}x] \implies x = y.$$

3.  $\mathcal{R}$  is *transitive* :

$$\forall x, y, z \in E, [x\mathcal{R}y \wedge y\mathcal{R}z] \implies x\mathcal{R}z.$$

**Definition 3.4.** Let  $\mathcal{R}$  be an order on  $E$ .

- An order relation  $\mathcal{R}$  on a set  $E$  is total if:

$$\forall x, y \in E : x\mathcal{R}y \text{ ou } y\mathcal{R}x.$$

It is also called  $(E, \mathcal{R})$  a totally ordered set.

- If the order  $\mathcal{R}$  is not total, we say that  $\mathcal{R}$  is a partial order.

**Example 3.3.** We equip  $\mathbb{R}^2$  with the relation noted as  $\mathcal{R}$  defined by:

$$(x, y)\mathcal{R}(x', y') \implies x \leq x' \text{ et } y \leq y'.$$

Demonstrate that  $\mathcal{R}$  is order relation on  $\mathbb{R}^2$ . Is the order total?

1.  $\mathcal{R}$  is *reflexive* :

$$\text{Let } (x, y) \in \mathbb{R}^2, \text{ we have } x \leq x \text{ and } y \leq y \implies (x, y)\mathcal{R}(x, y) .$$

2.  $\mathcal{R}$  is *anti symmetric* :

Let  $(x, y), (x', y') \in \mathbb{R}^2$ , we have  $(x, y)\mathcal{R}(x', y')$  and  $(x', y')\mathcal{R}(x, y)$ , then we have both  $x \leq x'$  and  $x' \leq x$  then  $x = x'$  and likewise  $y = y'$ .

3.  $\mathcal{R}$  is *transitive* :

Let  $(x, y), (x', y'), (x'', y'') \in \mathbb{R}^2$ , we have  $(x, y)\mathcal{R}(x', y')$  and  $(x', y')\mathcal{R}(x'', y'')$ , then we have both  $x \leq x' \leq x''$  and  $y \leq y' \leq y''$  then  $(x, y)\mathcal{R}(x'', y'')$ .

So  $\mathcal{R}$  is order relation .

The order is not total, because we cannot compare  $(0, 1)$  and  $(1, 0)$ .